

"Formulações estendidas para o problema  
do caminho mais curto com  
restrições de salto"

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Ciência 2010

## Hard to find if Problem is NP-Hard

$$\begin{array}{l} \min cx \\ \text{s.t. } Ax \geq b \\ x \geq 0 \text{ e inteiro} \end{array}$$

$$\begin{array}{l} \min cx \\ \text{s.t. } \text{Conv}\{x : Ax \geq b \\ x \geq 0 \text{ e inteiro}\} \end{array}$$

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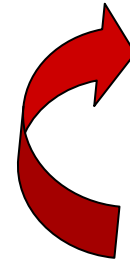
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**Projection**

$$\begin{array}{l} Bx + Cw \geq d \\ x, w \geq 0 \end{array}$$

# Path Reformulations

$$\min cx$$

$$s.a \quad A1x \geq b1$$

$$A2x \geq b2$$

$$x \geq 0 \quad e \quad \text{int } eiro$$

$$A2x \geq b2$$

$$x \geq 0 \quad e \quad \text{int } eiro$$



Path in an Expanded Graph

Usual Exact Formulation in Expanded Graph

Expanded Graph of Polynomial Size

Extended Exact Formulation in Original Graph

Compact Formulation

$$Bx + Cw \geq d$$

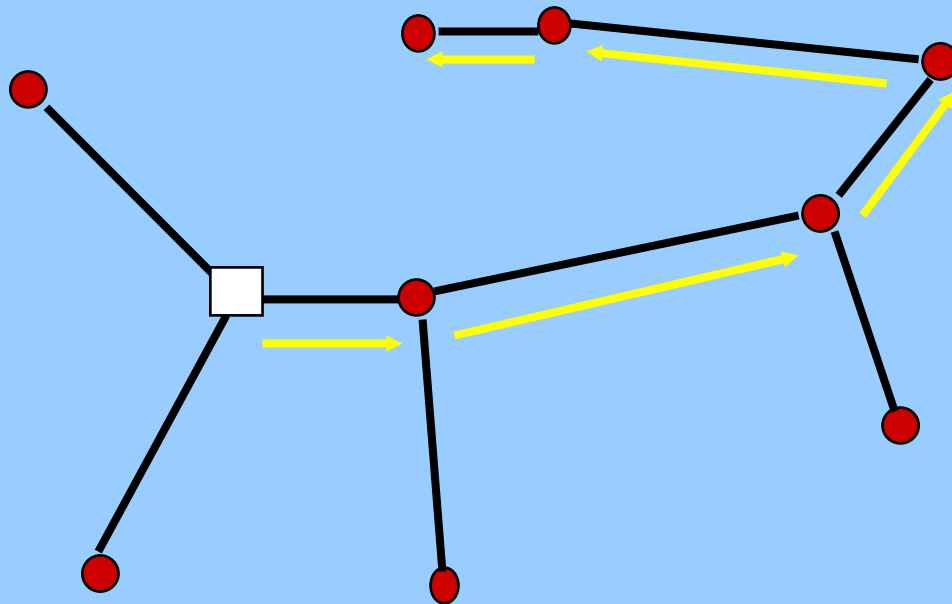
$$x, w \geq 0$$

## Hop-Constrained Spanning Tree =

Spanning Tree +

for each  $j \neq 0$ , n° of arcs in the path from node 0 to node  $j$  is  $\leq H$

Unfeasible solution if  $H < 5$



NP-Hard

Special case of  
Location problem  
( $H=2$ )

# Hop Constrained Spanning Trees

*min*

$$\sum_{i=0}^n \sum_{j=1}^n c_{ij} x_{ij}$$

*s. a.*

$$\sum_{i=0}^n x_{ij} = 1 \quad j = 1, \dots, n$$

$\{(i,j): i=0, \dots, n; j = 1, \dots, n; y_{ij}^k = 1\}$   
contains a path from 0 to k with at  
most H hops  $\forall k$

“contains” ?!

$$y_{ij}^k \leq x_{ij} \quad i = 0, \dots, n; j = 1, \dots, n$$

$$x_{ij} \in \{0,1\} \quad i = 0, \dots, n; j = 1, \dots, n$$



## HC-Path(k)

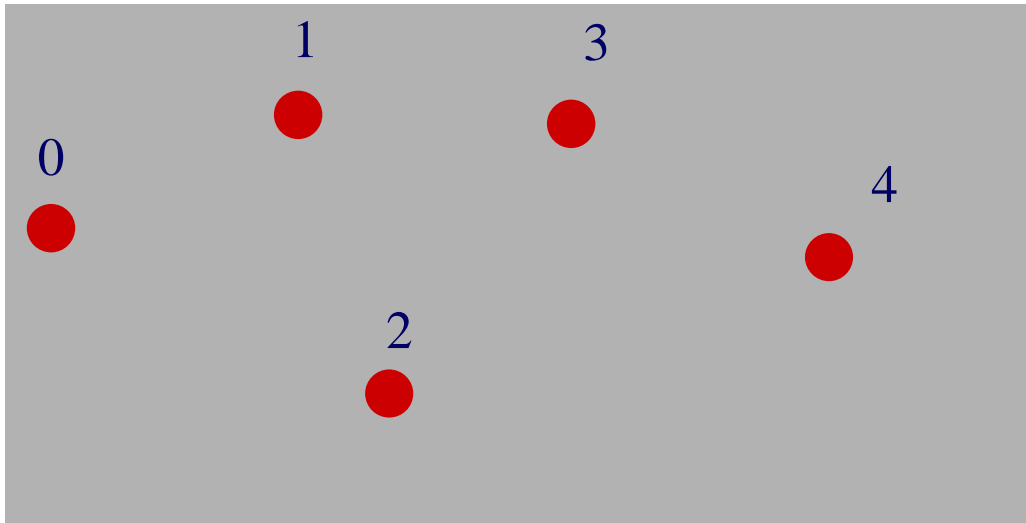
$$\sum_{i=1}^n f_{ij}^k - \sum_{i=1}^n f_{ji}^k = \begin{cases} -1 & \text{se } j=1 \\ 0 & \text{se } j=2, \dots, n; j \neq k \\ 1 & \text{se } j=k \end{cases}$$

$$\sum_{i=0}^n \sum_{j=1}^n f_{ij}^k \leq H \quad !?$$

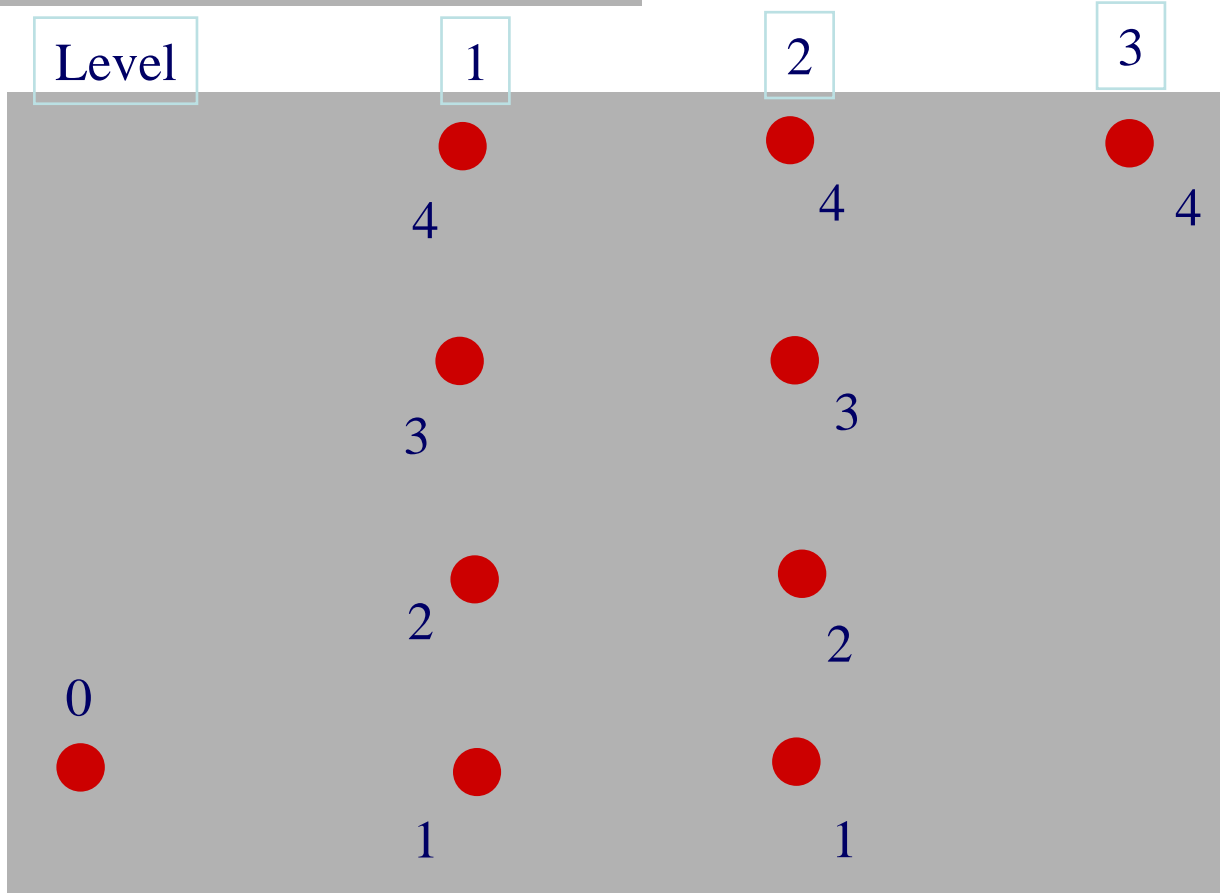
$$f_{ij}^k \in \{0,1\} \quad i = 0, \dots, n; j = 1, \dots, n$$

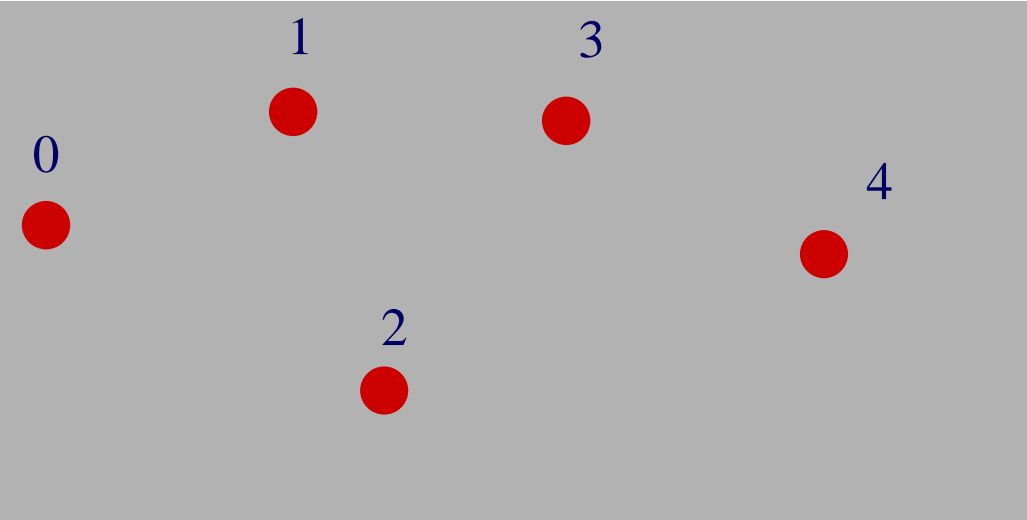
**In general,  $F(\text{HC-PATH}(k)_L)$   
contains fractional extreme points**

“Exactly” modelled in  
an adequate expanded graph (of polynomial size)

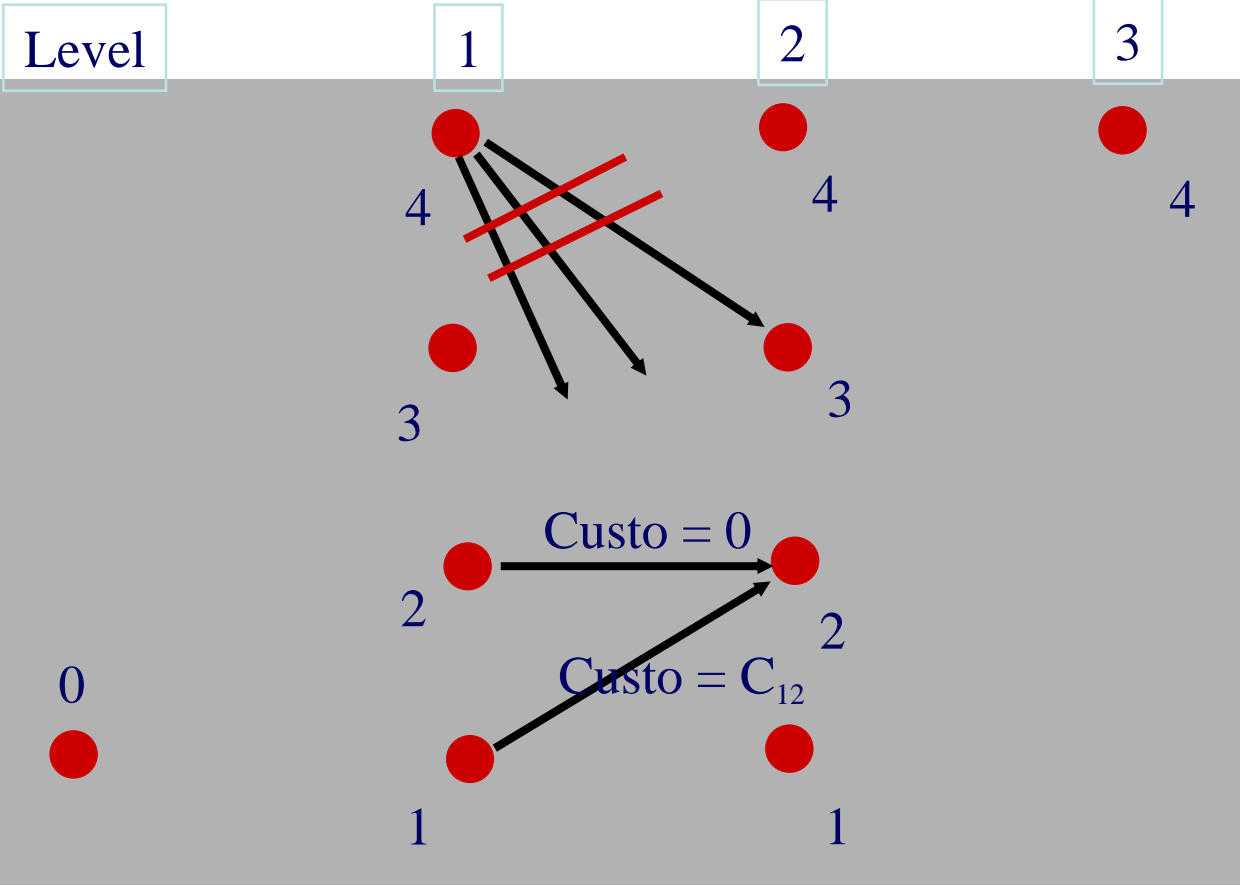


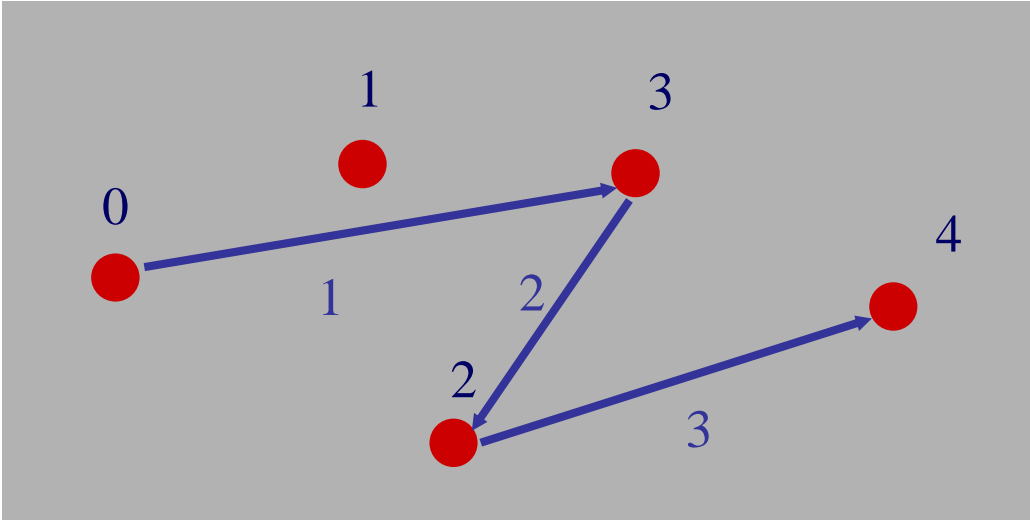
H=3



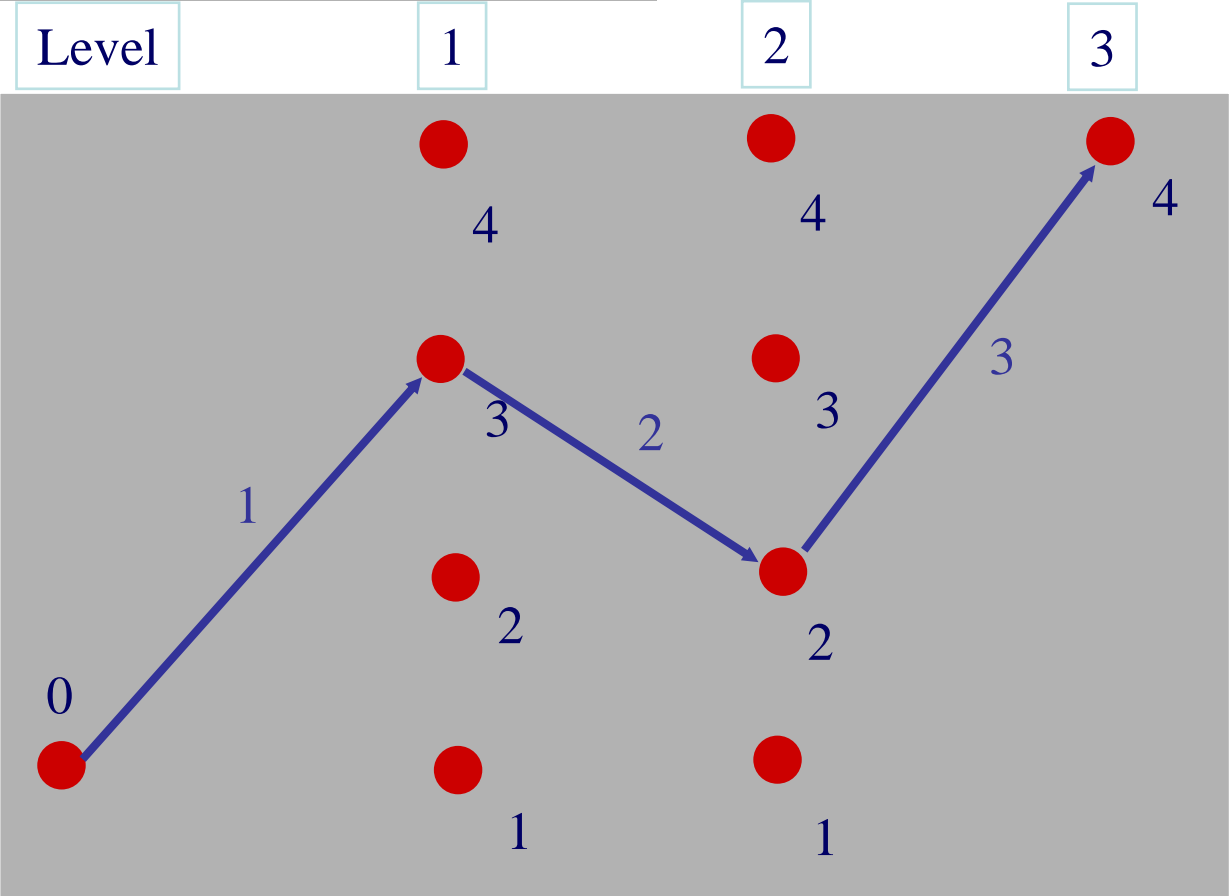


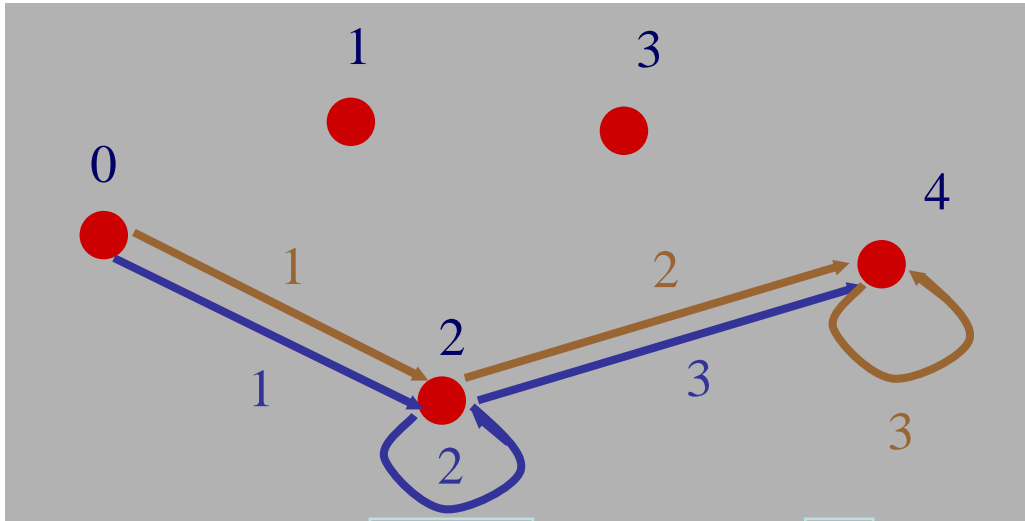
H=3





H=3





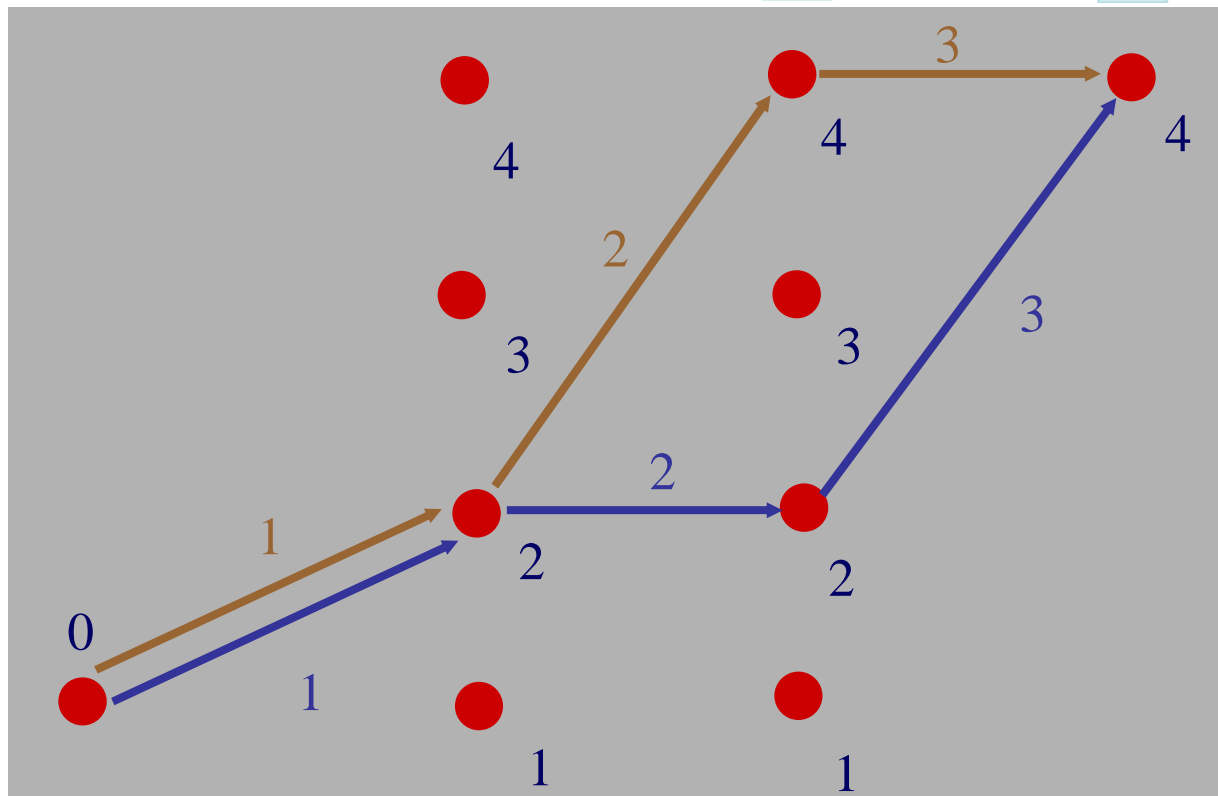
H=3

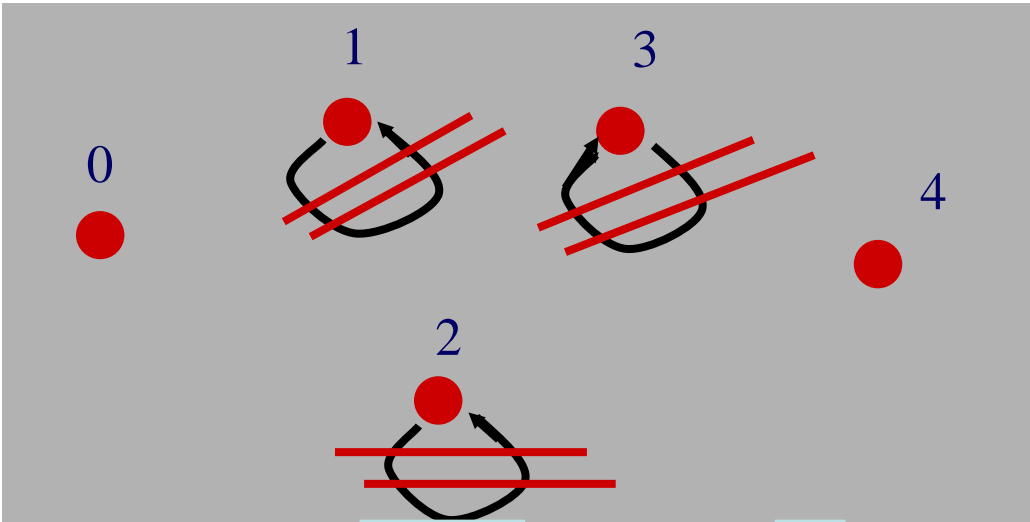
Level

1

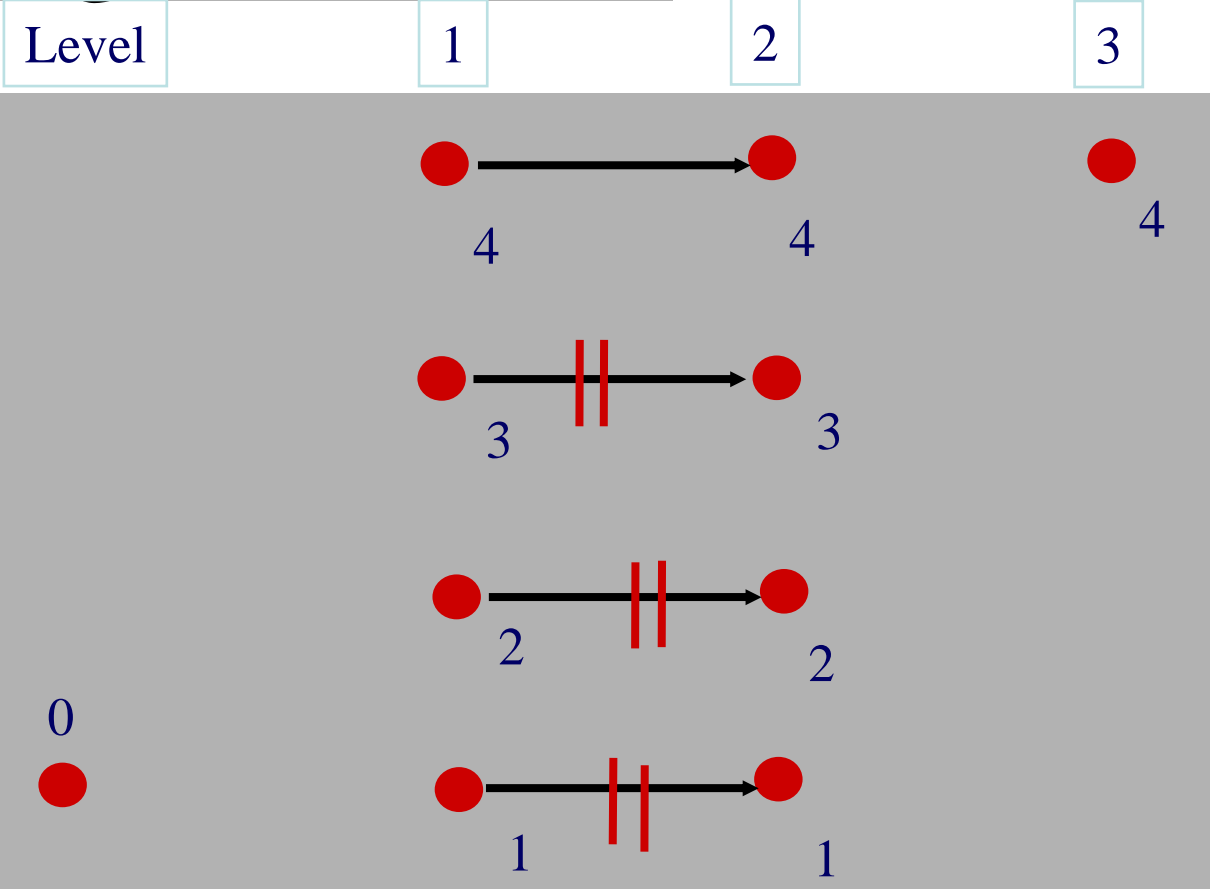
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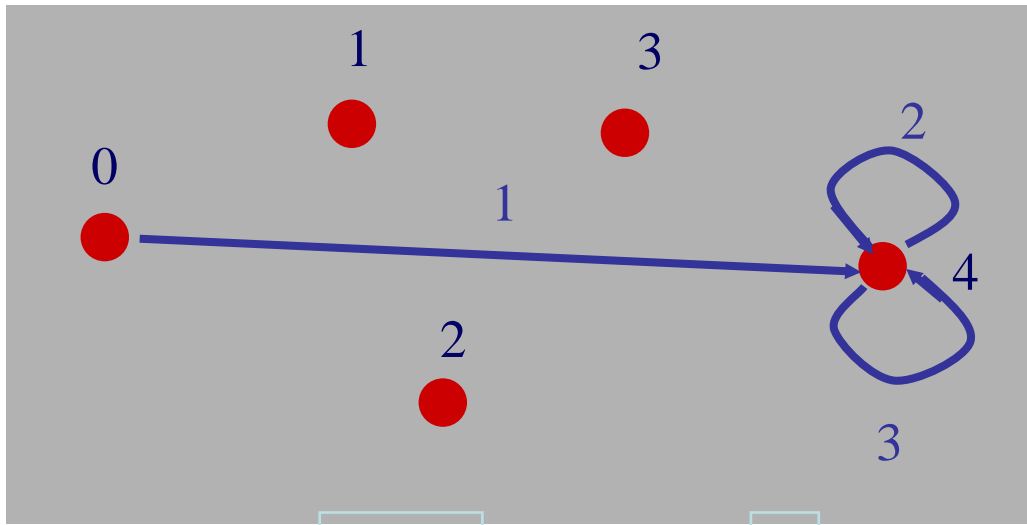
3





H=3



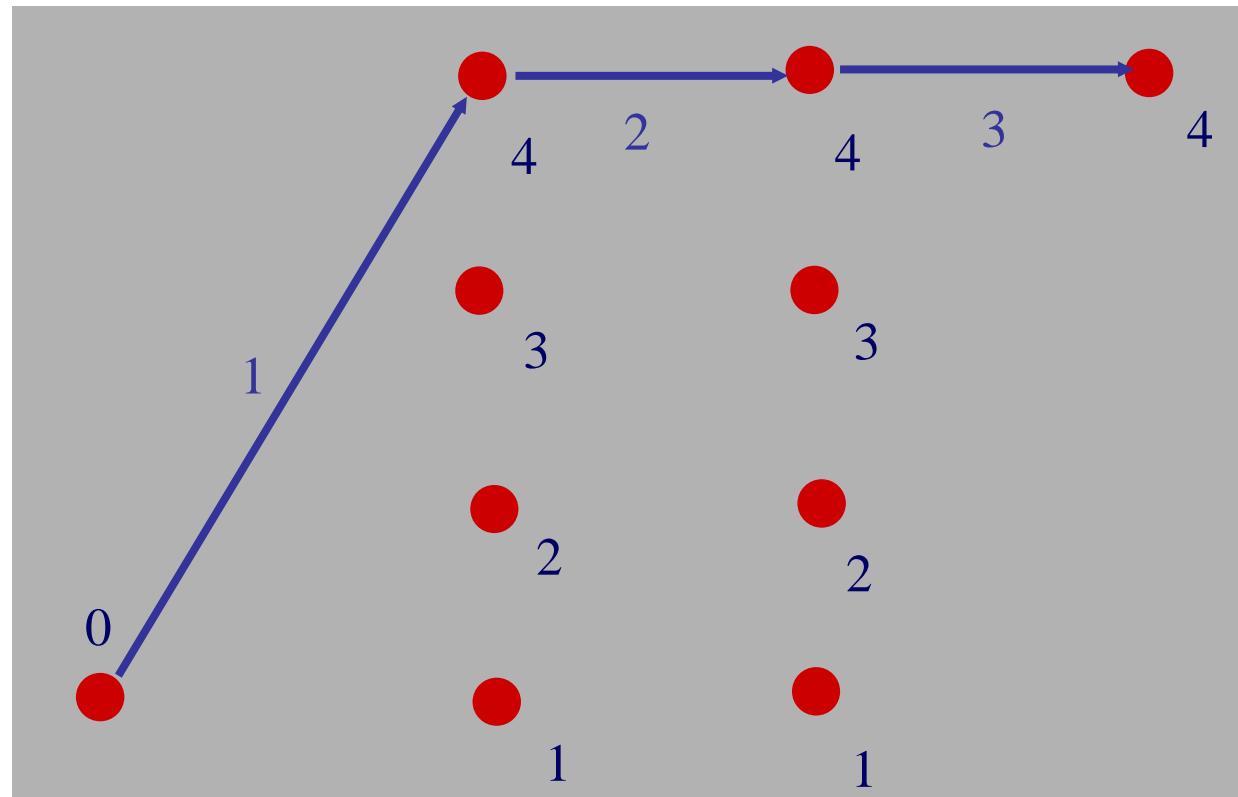


Level

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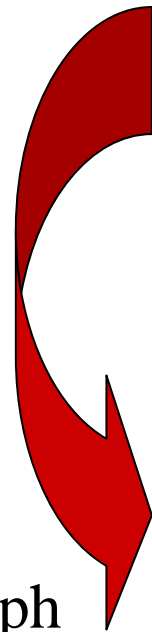
2

3



Expanded Graph

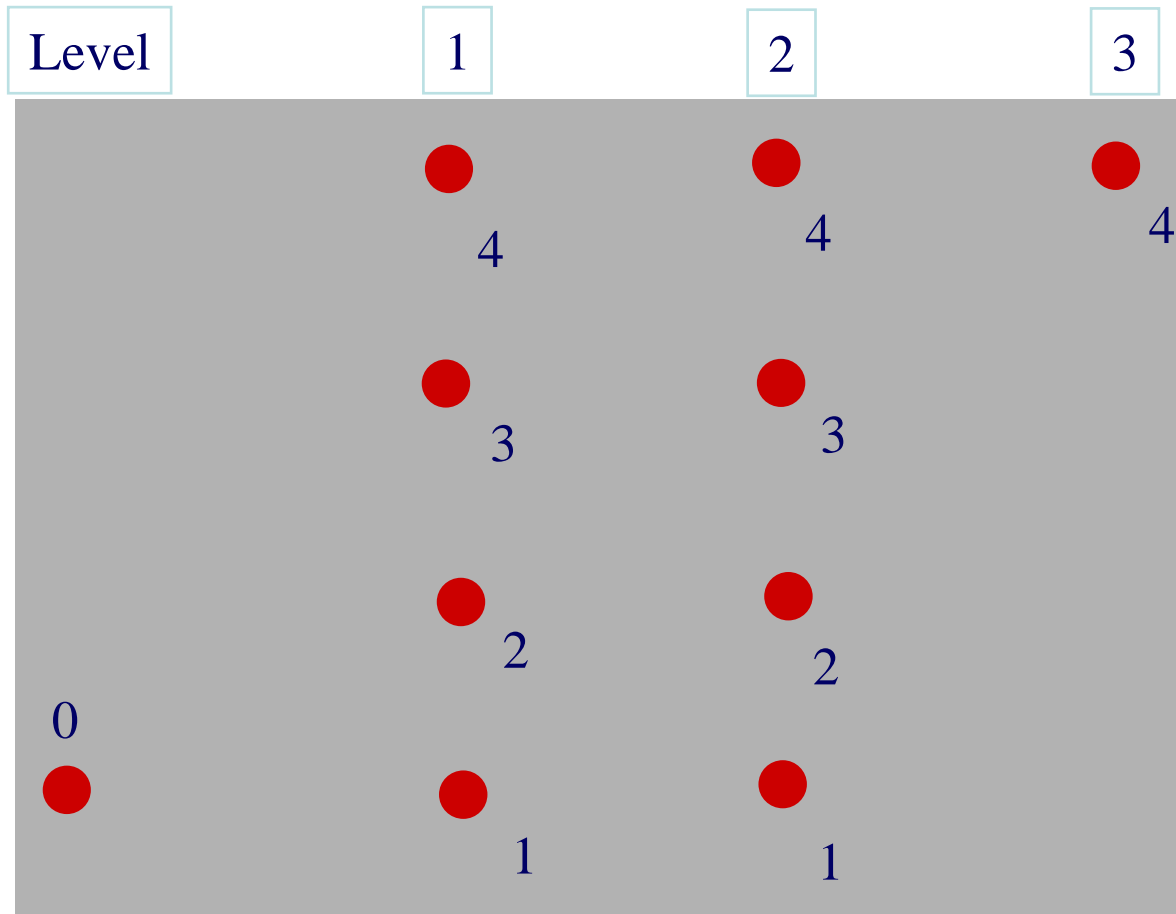
Unconstrained Path:  
Usual Exact Formulation in Expanded Graph  
Expanded Graph of Polynomial Size



Original Graph

Hop-Constrained Path  
Extended Exact Formulation in Original Graph  
Compact (of polynomial size) Formulation





Expanded Graph  $z_{i^{h-1}j^h} \in \{0,1\}$  indicating whether arc  $(i^{h-1},j^h)$  is in the solution

Original Graph  $z_{ij}^h \in \{0,1\}$  indicating whether arc  $(i,j)$  is in position  $h$  in the solution

# EXACT-HC-Path(k)

$$\min \sum_{i=0, \dots, n} \sum_{j=1, \dots, n} \sum_{h=1, \dots, H} d_{ij} Z_{ij}^{hk}$$

$$\begin{aligned} d_{ij} &= C_{ij} & i \neq j \\ d_{nn} &= 0 \end{aligned}$$

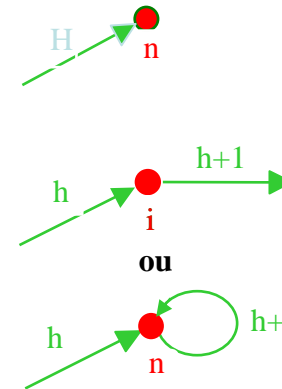
$$s.a \quad \sum_{j=1, \dots, n} Z_{jk}^{Hk} = 1$$

$$\sum_{j=1, \dots, n} Z_{ji}^{hk} = \sum_{j=1, \dots, n} Z_{ij}^{h+1, k} \quad i = 1, \dots, n;$$

$$Z_{0i}^{1k} = \sum_{j=1, \dots, n} Z_{ij}^{2k} \quad i = 1, \dots, n$$

$$Z_{ij}^{qk} \in \{0, 1\} \quad i = 0, \dots, n; j = 1, \dots, n, q = 1, \dots, H$$

$$\sum_{q=1, \dots, H} Z_{ij}^{qk} = Y_{ij}^k \quad i = 0, \dots, n; j = 1, \dots, n$$



MCF – generic with HC-Path(k) for each k

HDMCF - generic with Exact-HC-Path(k) for each k

Root	n	m	H	$v(\text{MCF}_L)$	$v(\text{HDMCF}_L)$	OPT
TC	60	210	4	776.5 (3868)	787.3 (21)	788 (5)
TE	60	210	4	1329.9 (9962)	1404.5 (125)	1407 (110)